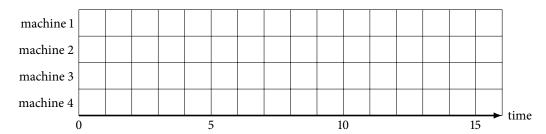
## **Lesson 9. Machine Scheduling**

**Example 1.** The Markov Micromanufacturing Company has 9 production jobs it needs to process in the next 24 hours. The company has 4 identical machines that run in parallel. Each of these 9 jobs must be run on one of these machines **nonpreemptively**: that is, once a job is started on a machine, it must stay on that machine until it is completed. The processing times of these jobs are given below:

job	1	2	3	4	5	6	7	8	9
processing time (hours)	7	7	6	6	5	5	4	4	4

The company wants to minimize the makespan, or the completion time of the last job to finish processing.

- Let m be the number of machines in this case, m = 4
- Suppose we schedule the jobs using the **longest processing time first (LPT)** rule:
  - First, schedule the *m* longest jobs on the *m* machines
  - o Whenever a machine becomes free, put the longest unprocessed job on that machine
- Idea: LPT puts shorter jobs towards the end of the schedule, where they can be used to balance the loads on each machine
- For our problem, this yields a schedule that looks like this:



- o This kind of diagram is known as a Gantt chart
- Therefore, the makespan for the LPT schedule is
- It turns out that the makespan of an LPT schedule is always at most  $33.\overline{3}\%$  larger than the minimum makespan
- So... can we do better?
- Let's formulate this problem as a dynamic program
- Stages:

Decisions, transitions, and rewards/costs at stage t (edges):  Shortest/longest path?  Winimum makespan ↔				
Shortest/longest path?				
	Decisions, transitions, an	d rewards/costs at sta	ge t (edges):	
Minimum makespan ←→	11 ( (/1 ) ( 1.3			
Minimum makespan ↔	nortest/longest path?			
Assignments of jobs to machines ←→				
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## A Problems

See the accompanying Jupyter Notebook for this lesson.